

**USEFUL FOR UNIVERSITY EXAMS, GATE,
NET AND OTHER CS EXAMS**

DATABASE MANAGEMENT SYSTEM

**RELATIONAL
ALGEBRA with
EXAMPLE**

**PROPER
NOTES
IN PPT
FORM**

**PART -
10**



Relational Algebra

- ▶ Relational algebra is a procedural query language.
- ▶ In this we have to give a step by step process to obtain the result of the query.
- ▶ It uses operators to perform queries.

1. Select Operation:

- ▶ The select operation selects tuples that satisfy a given predicate(condition).
- ▶ It is denoted by sigma (σ).
- ▶ Notation: $\sigma p(r)$

Here :

σ is used for selection operator for tuples.

r is used for relation

p is used as a propositional logic formula which may use connectors like: AND OR and NOT. These relational can use as relational operators like $=, \neq, \geq, <, >, \leq$.

EXAMPLE

- ▶ Consider the following relational database schema consisting of the four relation schemas:

passenger (pid, pname, pgender, pcity)

agency (aid, aname, acity)

flight (fid, fdate, time, src, dest)

booking (pid, aid, fid, fdate)

- ▶ **Get the complete details of all flights to New Delhi.**

$\sigma_{\text{destination} = \text{"New Delhi"}}(\text{flight})$

Project Operation:

- ▶ This operation shows the list of those attributes that we wish to appear in the result. Rest of the attributes are eliminated from the table.
- ▶ It is denoted by Π .
- ▶ Notation: $\Pi A_1, A_2, A_n (r)$
- ▶ **Where**
A1, A2, A3 is used as an attribute name of relation **r**.
- ▶ EXAMPLE : $\Pi \text{NAME, CITY (EMPLOYER)}$

Union Operation:

- ▶ Suppose there are two relations R and S.
- ▶ The union operation contains all the tuples that are either in R or S or both in R & S.
- ▶ It eliminates the duplicate tuples. It is denoted by \cup .
- ▶ R and S must have the attribute of the same number.
- ▶ Duplicate tuples are eliminated automatically.
- ▶ EXAMPLE :

Π CUSTOMER_NAME (BORROW) \cup Π CUSTOMER_NAME (DEPOSITOR)

Set Intersection:

- ▶ Suppose there are two relations R and S.
- ▶ The set intersection operation contains all tuples that are in both R & S.
- ▶ It is denoted by intersection \cap .
- ▶ Notation: $R \cap S$

EXAMPLE:

Π CUSTOMER_NAME (BORROW) \cap Π CUSTOMER_NAME (DEPOSITOR)

Set Difference:

- ▶ Suppose there are two relations R and S.
- ▶ The set intersection operation contains all tuples that are in R but not in S.
- ▶ It is denoted by intersection minus (-).
- ▶ Notation: $R - S$

EXAMPLE:

$\Pi \text{ CUSTOMER_NAME (BORROW)} - \Pi \text{ CUSTOMER_NAME (DEPOSITOR)}$

Cartesian product

- ▶ The Cartesian product is used to combine each row in one table with each row in the other table.
- ▶ It is also known as a cross product.
- ▶ It is denoted by \times .
- ▶ Notation: $E \times S$

Example

EMP_CODE	EMP_NAME
101	A
102	B
103	C

EMP_CODE	SALARY
101	1000
102	2000
103	3000

e.EMP_CODE	e.EMP_NAME	s.Emp_CODE	s.SALARY
101	A	101	1000
101	A	102	2000
101	A	103	3000
102	B	101	1000
102	B	102	2000
102	B	103	3000
103	C	101	1000
103	C	102	2000
103	C	103	3000

Join Operations:

- ▶ A Join operation combines related tuples from different relations, if and only if a given join condition is satisfied.
- ▶ It is denoted by \bowtie .

Example

EMP_CODE	EMP_NAME
101	A
102	B
103	C

EMP_CODE	SALARY
101	1000
102	2000
103	3000

e.EMP_CODE	e.EMP_NAME	s.Emp_CODE	s.SALARY
101	A	101	1000
101	A	102	2000
101	A	103	3000
102	B	101	1000
102	B	102	2000
102	B	103	3000
103	C	101	1000
103	C	102	2000
103	C	103	3000

EXAMPLES

- Consider the following relational database schema consisting of the four relation schemas:

passenger (pid, pname, pgender, pcity)

agency (aid, aname, acity)

flight (fid, fdate, time, src, dest)

booking (pid, aid, fid, fdate)

- Get the details about all flights from Chennai to New Delhi.

$\sigma_{src = \text{"Chennai"} \wedge dest = \text{"New Delhi"}}(\text{flight})$

EXAMPLES

- Consider the following relational database schema consisting of the four relation schemas:

passenger (pid, pname, pgender, pcity)

agency (aid, aname, acity)

flight (fid, fdate, time, src, dest)

booking (pid, aid, fid, fdate)

- Find only the flight numbers for passenger with pid 123 for flights to Chennai before 06/11/2020.

$\Pi_{fid} (\sigma_{pid = 123} (\text{booking}) \bowtie \sigma_{dest = \text{"Chennai"} \wedge fdate < 06/11/2020} (\text{flight}))$

Given conditions are pid, dest, and fdate. To get the flight id for a passenger given a pid, we have two tables flight and booking to be joined with necessary conditions. From the result, the flight id can be projected]

EXAMPLES

- Consider the following relational database schema consisting of the four relation schemas:

passenger (pid, pname, pgender, pcity)

agency (aid, aname, acity)

flight (fid, fdate, time, src, dest)

booking (pid, aid, fid, fdate)

- Find the **passenger names** for passengers who have bookings on at least one flight.

$\Pi_{pname}(\text{passenger} \bowtie \text{booking})$

EXAMPLES

- Consider the following relational database schema consisting of the four relation schemas:

passenger (pid, pname, pgender, pcity)

agency (aid, aname, acity)

flight (fid, fdate, time, src, dest)

booking (pid, aid, fid, fdate)

- Find the **passenger names** for those who do not have any bookings in any flights.

$\Pi_{pname} ((\Pi_{pid}(\text{passenger}) - \Pi_{pid}(\text{booking})) \bowtie \text{passenger})$

here applied a set difference operation. The set difference operation returns only pids that have no booking. The result is joined with passenger table to get the passenger names

EXAMPLES

- Consider the following relational database schema consisting of the four relation schemas:

passenger (pid, pname, pgender, pcity)

agency (aid, aname, acity)

flight (fid, fdate, time, src, dest)

booking (pid, aid, fid, fdate)

- Find the **agency names** for agencies that located in the same city as passenger with **passenger id 123**.

$\Pi_{aname} (agency \bowtie_{acity = pcity} (\sigma_{pid = 123} (passenger)))$

THANK YOU

